## Research Article

# On chromatic vertex stability of 3-chromatic graphs with maximum degree 4 

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#### Abstract

The (independent) chromatic vertex stability (ivs $(G)) \operatorname{vs}_{\chi}(G)$ is the minimum size of (independent) set $S \subseteq V(G)$ such that $\chi(G-S)=\chi(G)-1$. The question of how large must the chromatic number $\chi(G)$ of a graph $G$ be, in terms of the maximum degree $\Delta(G)$, to ensure the equality $\operatorname{ivs}_{\chi}(G)=\mathrm{vs}_{\chi}(G)$ was raised by Akbari et al. [European J. Combin. 102 (2022) \#103504]; the authors showed that $\operatorname{ivs}_{\chi}(G)=\mathrm{vs}_{\chi}(G)$ if $\chi(G) \in\{\Delta(G), \Delta(G)+1\}$, and also pointed out to graphs with $\chi(G) \leq(\Delta(G)+1) / 2$ for which $\operatorname{ivs}_{\chi}(G)>\mathrm{vs}_{\chi}(G)$. In the light of their findings, they raised the following problem: Is it true that $\chi(G) \geq \Delta(G) / 2+1$ always implies ivs $\chi_{\chi}(G)=\mathrm{vs}_{\chi}(G)$ ? This threshold question was recently answered in the negative by Cambrie et al. [arXiv: 2203.13833 v 1 , (2022)]. In this paper, we show that the smallest instance for counterexamples is the case $(\chi(G), \Delta(G))=(3,4)$, with the smallest possible order being 9 (and there are 30 such graphs). We construct exponentially many graphs $G$ having $\Delta(G)=4, \chi(G)=3, \operatorname{ivs}_{\chi}(G)=3$, and vs $\chi(G)=2$.


Keywords: chromatic vertex stability; independent chromatic vertex stability; chromatic number; maximum degree.
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## 1. Introduction

Let $G$ be a graph. Its edge stability number, $\mathrm{es}_{\chi}(\mathrm{G})$, is the minimum number of edges whose deletion results in a graph $H$ with $\chi(H)=\chi(G)-1$. The edge stability number was introduced in 1980 by Staton [9], and rediscovered in 2008 by Arumugam, Hamid, and Muthukamatchi [4]. For recent results on this invariant, see e.g. [1, 3, 6, 8].

General concept of stability number appeared in [5], but the first paper on chromatic vertex stability number was written by Akbari, Beikmohammadi, Klavžar, and Movarraei in 2021, see [2]. The chromatic vertex stability vs ${ }_{\chi}(G)$ of $G$ is the minimum number of vertices of $G$ such that their deletion results in a graph $H$ with $\chi(H)=\chi(G)-1$. Analogously, the independent chromatic vertex stability $\operatorname{ivs}_{\chi}(G)$ of $G$ is the minimum number of independent vertices of $G$ such that their deletion results in a graph $H$ with $\chi(H)=\chi(G)-1$. Obviously, $\operatorname{vs}_{\chi}(G) \leq \operatorname{ivs}_{\chi}(G)$. The main result of [2] is the following.

Theorem 1.1. If $G$ is a graph with $\chi(G) \in\{\Delta(G), \Delta(G)+1\}$ then $\operatorname{vs}_{\chi}(G)=\operatorname{ivs}_{\chi}(G)$.
The authors defined the threshold function $f(\Delta)$ as the smallest quantity such that, for any graph $G$ of maximum degree $\Delta$, it must hold that $\mathrm{vs}_{\chi}(G)=\operatorname{ivs}_{\chi}(G)$ provided that $\chi(G) \geq f(\Delta)$. Notice that Theorem 1.1 asserts $f(\Delta) \leq \Delta$. They also showed that as soon as $\chi(G) \leq(\Delta(G)+1) / 2$ the equality $\operatorname{vs}_{\chi}(G)=\operatorname{ivs}_{\chi}(G)$ need no longer be true, and consequently asked the following question (see Problem 3.2 in [2]).

Problem 1.1. Is it true that $f(\Delta) \leq \frac{\Delta(G)}{2}+1$, that is, if $G$ is a graph with $\chi(G) \geq \frac{\Delta(G)}{2}+1$ does it then always hold that $\operatorname{vs}_{\chi}(G)=\operatorname{ivs}_{\chi}(G)$ ?

The question of Problem 1.1 was recently answered in the negative by Cambie et al. [7], who proved that $f(\Delta)=\Delta$ for $3 \leq \Delta \leq 10$. Additionally, they determined the threshold $f(\Delta)$ to within two values (and indeed sometimes a unique value) for graphs of sufficiently large maximum degree.

In this paper, we focus on the smallest value $\chi(G)=3$ for which counterexamples to Problem 1.1 exist, i.e. we consider the case when $\chi(G)=3$ and $\Delta(G)=4$. A simple 'ladder-like' counterexample on 9 vertices is depicted in Figure 1. Notice that the ladder part of the counterexample can be of any length $4 k$ for $k \geq 2$.

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Figure 1: A graph $G$ with $\Delta(G)=4$ and $\chi(G)=3$. The blue independent 3 -set realizes ivs $(G)=3$ and the yellow 2 -set realizes $\mathrm{vs}_{\chi}(G)=2$.

Our main result is the following.
Theorem 1.2. For each $n \geq 9$ there are at least $\max \left\{1,2^{\left\lfloor\frac{n-11}{2}\right\rfloor}\right\}$ planar graphs with $\chi(G)=3, \Delta(G)=4$, $\mathrm{ivs}_{\chi}(G)=3$ and $\mathrm{vs}_{\chi}(G)=2$.

Observe that if $\chi(G)=3$ and $\Delta(G)=4$, then $\chi(G)=\frac{\Delta(G)}{2}+1=\Delta(G)-1$. Hence, the bound on $\chi(G)$ in Theorem 1.1 cannot be relaxed when $\chi(G)=3$.

## 2. Proofs

We start with a pair of simple observations followed by a couple of lemmas.
Observation 2.1. For every graph $G$, ivs ${\underset{\chi}{\chi}}(G)$ equals the minimum size of a colour class over all proper $\chi(G)$-colourings of $G$. Hence $|V(G)| \geq \operatorname{ivs}_{\chi}(G) \cdot \chi(G)$.

Proof. Let us first notice that there exists a proper $\chi(G)$-colouring of $G$ with a colour class of size ivs $\chi_{\chi}(G)$. Indeed, take $S \subseteq V(G)$ to be an independent set of size $|S|=\operatorname{ivs}_{\chi}(G)$ such that $\chi(G-S)=\chi(G)-1$. Use a proper $(\chi(G)-1)$-colouring of $G-S$ and assign to all vertices of $S$ a new colour. So ivs $_{\chi}(G)$ is not less than the minimum size of a colour class over all proper $\chi(G)$-colourings of $G$.

Contrarily, consider a proper $\chi(G)$-colouring of $G$ which minimizes the size of a colour class, and let $S$ be such a minimum colour class. Then $S$ is an independent subset of $V(G)$ and $\chi(G-S) \leq \chi(G)-1$. In fact, we must have equality here for otherwise $G$ would admit a proper $(\chi(G)-1)$-colouring. So ivs $_{\chi}(G)$ is also not more than the minimum size of a colour class over all proper $\chi(G)$-colourings of $G$, which proves our point.

The inequality $|V(G)| \geq \operatorname{ivs}_{\chi}(G) \cdot \chi(G)$ is now an immediate consequence.
Observation 2.2. If $\Delta(G) \leq 2$ then $\operatorname{vs}_{\chi}(G)=\operatorname{ivs} \chi_{\chi}(G)$.
Proof. We may assume that $\mathrm{vs}_{\chi}(G) \geq 2$. Indeed, if $\mathrm{vs}_{\chi}(G)=1$ then obviously $\mathrm{ivs}_{\chi}(G)=1$ as well. We may also assume that $G$ is connected. Then $G$ is either a path or an even cycle. In either case

$$
\operatorname{vs}_{\chi}(G)=\operatorname{ivs}_{\chi}(G)=\left\lfloor\frac{|V(G)|}{2}\right\rfloor .
$$

As already mentioned, we are interested in finding graphs $G$ for which $\chi(G) \geq \frac{\Delta(G)}{2}+1$ and ivs $\chi_{\chi}(G)>\mathrm{vs}_{\chi}(G)$. Our first lemma establishes some implications for the order and the considered stability parameters.

Lemma 2.1. If $\mathrm{ivs}_{\chi}(G)>\mathrm{vs}_{\chi}(G)$ and $\chi(G) \geq \frac{\Delta(G)}{2}+1$ then $|V(G)| \geq 9, \mathrm{ivs}_{\chi}(G) \geq 3, \mathrm{vs}_{\chi}(G) \geq 2$ and $\chi(G) \geq 3$. Moreover, if $|V(G)|=9$ then $\operatorname{ivs}_{\chi}(G)=3$, vs $_{\chi}(G)=2$ and $\chi(G)=3$.

Proof. Since ivs $\chi_{\chi}(G)>\operatorname{vs}_{\chi}(G)$, we must have $\mathrm{vs}_{\chi}(G) \geq 2$ and consequently ivs $(G) \geq 3$.
If $\chi(G) \leq 2$ then from $\chi(G) \geq \frac{\Delta(G)}{2}+1$ we get $\Delta(G) \leq 2(\chi(G)-1) \leq 2$, which in view of Observation 2.2 contradicts $\operatorname{ivs}_{\chi}(G)>\mathrm{vs}_{\chi}(G)$. Hence $\chi(G) \geq 3$.

From the inequality stated in Observation 2.1, it follows that $|V(G)| \geq \operatorname{ivs}_{\chi}(G) \cdot \chi(G) \geq 3 \cdot 3$, that is, $|V(G)| \geq 9$. And if $|V(G)|=9$ then $\chi(G)=\operatorname{ivs}_{\chi}(G)=3$ and $\mathrm{vs}_{\chi}(G)=2$.

Figure 2 depicts two graphs, respectively denoted by $G_{9}$ and $G_{10}$ in regard to their orders. The former one can be obtained from the octahedron graph by subdividing the edges of a triangle. It has $\Delta\left(G_{9}\right)=4, \chi\left(G_{9}\right)=\operatorname{ivs}_{\chi}\left(G_{9}\right)=3$ and $\operatorname{vs}_{\chi}\left(G_{9}\right)=2$. Observe that $\chi\left(G_{9}-\{x, y\}\right)=2$ if and only if $\{x, y\}=\left\{v_{i}, v_{j}\right\}$, where $1 \leq i<j \leq 3$, and for an independent set of vertices $\{x, y, z\}$ we have $\chi\left(G_{9}-\{x, y, z\}\right)=2$ if and only if $\{x, y, z\}=\left\{u_{i}, v_{i}, w_{i}\right\}$, where $1 \leq i \leq 3$. The graph $G_{10}$ is obtained from $G_{9}$ by adding the vertex $q$ and connecting it to $w_{2}$ and $w_{3}$. It also has $\Delta\left(G_{10}\right)=4, \chi\left(G_{10}\right)=\operatorname{ivs} \chi\left(G_{10}\right)=3$ and $\operatorname{vs}_{\chi}\left(G_{10}\right)=2$. Again $\chi\left(G_{10}-\{x, y\}\right)=2$ if and only if $\{x, y\}=\left\{v_{i}, v_{j}\right\}$, where $1 \leq i<j \leq 3$, and for an independent set of vertices $\{x, y, z\}$ we have $\chi\left(G_{10}-\{x, y, z\}\right)=2$ if and only if $\{x, y, z\}=\left\{u_{i}, v_{i}, w_{i}\right\}$ where $2 \leq i \leq 3$.


Figure 2: The graph $G_{9}$ (left) and the graph $G_{10}$ (right).
Our second lemma concerns the case $\chi(G)=\operatorname{ivs}_{\chi}(G)=3$ and $\mathrm{vs}_{\chi}(G)=2$. Under the assumption $\chi(G) \geq \frac{\Delta(G)}{2}+1$ we establish the maximum degree of $G$ and the vertex degrees of every 2-set which realizes vs $(G)$.

Lemma 2.2. Let $G$ be a graph with $\chi(G)=\operatorname{ivs}_{\chi}(G)=3, \operatorname{vs}_{\chi}(G)=2$ and $\chi(G) \geq \frac{\Delta(G)}{2}+1$. Then $\Delta(G)=4$ and for every $v_{1}, v_{2} \in V(G)$ such that $\chi\left(G-\left\{v_{1}, v_{2}\right\}\right)=2$ we have $\operatorname{deg}_{G}\left(v_{1}\right)=\operatorname{deg}_{G}\left(v_{2}\right)=4$.

Proof. Let $G$ satisfy the assumptions of the lemma and let $v_{1}, v_{2} \in V(G)$ be such that $\chi\left(G-\left\{v_{1}, v_{2}\right\}\right)=2$. Then

$$
\Delta(G) \leq 2(\chi(G)-1)=4
$$

For argument's sake, suppose that $\operatorname{deg}_{G}\left(v_{1}\right) \leq 3$. Since ivs ${ }_{\chi}(G)=3>2$, we have $v_{1} v_{2} \in E(G)$. Further, since $\chi\left(G-v_{2}\right)=$ 3, there must be an odd cycle in $G-v_{2}$; moreover, every such cycle passes through $v_{1}$ since $G-\left\{v_{1}, v_{2}\right\}$ is bipartite. Consequently $\operatorname{deg}_{G}\left(v_{1}\right)=3$.

Let $u_{1}, u_{2}$ be the neighbours of $v_{1}$ in $G-v_{2}$. Since every odd cycle in $G-v_{2}$ passes through both $u_{1}, u_{2}$, we conclude that $G-\left\{u_{i}, v_{2}\right\}$ is bipartite as well, $i \in\{1,2\}$. From $\operatorname{ivs}_{\chi}(G)=3>2$ it follows that $u_{1} v_{2}, u_{2} v_{2} \in E(G)$. Moreover, $u_{1} u_{2} \notin E(G)$, for otherwise $v_{1}, v_{2}, u_{1}, u_{2}$ induces a $K_{4}$, implying $\chi(G) \geq 4$.

Since $\operatorname{ivs}_{\chi}(G)=3$, there must be an odd cycle in $G-\left\{u_{1}, u_{2}\right\}$. This cycle cannot pass through $v_{1} \operatorname{since} \operatorname{deg}_{G-\left\{u_{1}, u_{2}\right\}}\left(v_{1}\right)=1$. If this cycle does not pass through $v_{2}$ as well then it is in $G-\left\{v_{1}, v_{2}\right\}$ which means that $\chi\left(G-\left\{v_{1}, v_{2}\right\}\right)=3$, a contradiction. Hence, there is an odd cycle passing through $v_{2}$ in $G-\left\{u_{1}, u_{2}, v_{1}\right\}$, which means that $\operatorname{deg}_{G}\left(v_{2}\right) \geq 5$. This contradiction settles the lemma.

A computer search shows that there are precisely 30 graphs $G$ of order 9 and having $\Delta(G)=4, \chi(G)=3$, ivs $\chi_{\chi}(G)=3$ and $\mathrm{vs}_{\chi}(G)=2$. Several of them (including $G_{9}$ ) are planar and four are obtained by adding an edge to another graph from the same collection.

For every $n \geq 9$ let $S_{n}$ be the set of graphs $G$ on $n$ vertices such that $\Delta(G)=4, \chi(G)=3$, ivs $\chi_{\chi}(G)=3$ and vs $(G)=2$. Thus $G_{9} \in S_{9}$ and $G_{10} \in S_{10}$. By $C_{\chi}(G)$ we denote the set of vertices $x \in V(G)$ such that there is some $y \in V(G)$ for which $\chi(G-\{x, y\})=2$; note that every such $y$ is a neighbour of $x$. For example, $C_{\chi}\left(G_{9}\right)=C_{\chi}\left(G_{10}\right)=\left\{v_{1}, v_{2}, v_{3}\right\}$. In view of our next result, for every $n \geq 9$ there is a planar graph in $S_{n}$ which is topologically equivalent to $G_{9}$ or $G_{10}$.

Proposition 2.1. Let $G \in S_{n}$ and $e_{1}, e_{2}, \ldots, e_{t} \in E(G) \backslash E\left(\left[C_{\chi}(G)\right]\right)$, i.e., each $e_{i}$ has at most one endvertex in $C_{\chi}(G)$. Let $n_{1}, n_{2}, \ldots, n_{t}$ be positive even integers. For every $i, 1 \leq i \leq t$, subdivide $e_{i}$ with $n_{i}$ new vertices, and denote the resulting graph by $H$. Then $\Delta(H)=4, \chi(H)=3, \operatorname{ivs}_{\chi}(H)=3$ and $\mathrm{vs}_{\chi}(H)=2$. In other words, $H \in S_{n+\left(n_{1}+\cdots+n_{t}\right)}$.
Proof. Obviously $\Delta(H)=\Delta(G)=4$. Since $\chi(G)=3$, the graph $G$ has an odd cycle. Since to any edge of this cycle we added an even number (possibly zero) of vertices, $H$ also has an odd cycle; thus $\chi(H) \geq 3$. Moreover, if $S \subseteq V(G)$ is such that $G-S$ is bipartite then $H-S$ is bipartite as well. Hence, $\chi(H)=3, \operatorname{ivs}_{\chi}(H) \leq 3$ and $\operatorname{vs}_{\chi}(H) \leq 2$.

If there is some $v \in V(H)$ such that $\chi(H-v)=2$, then $v \notin V(G)$. So $v$ is obtained by subdividing an edge, say $x y$, of $G$. However, as $H-v$ is bipartite, both $G-x$ and $G-y$ are bipartite, a contradiction. Hence, $\operatorname{vs}_{\chi}(H)=2$.

Finally, let us show that $\operatorname{ivs}_{\chi}(H)=3$. Supposing the opposite, there are $u, v \in V(H)$ such that $\chi(H-\{u, v\})=2$ and $u v \notin E(H)$. It cannot be that both $u$ and $v$ are in $V(G)$, because we did not subdivide edges connecting vertices of $C_{\chi}(G)$. So we may assume that $v$ is obtained by subdividing an edge $x y$ of $G$, where $y \notin C_{\chi}(G)$. Since $H-\{u, v\}$ is bipartite, so is $H-\{u, y\}$. But then $u$ cannot be a vertex of $G$ as well. Hence, $u$ is obtained by subdividing an edge $w z$ of $G$. As $H-\{u, y\}$ is bipartite, so is $H-\{z, y\}$. However, this contradicts the fact that $y \notin C_{\chi}(G)$.

From Proposition 2.1 we deduce that $S_{n} \neq \emptyset$ for every $n \geq 9$. Indeed, if $n$ is odd then take $G_{9}$, subdivide the edge $u_{2} w_{1}$ with $n-9$ new vertices and denote the resulting graph by $G_{n}$. Analogously if $n$ is even then take $G_{10}$, subdivide the edge $u_{2} w_{1}$ with $n-10$ new vertices and denote the resulting graph by $G_{n}$. Then $G_{n}$ is a connected planar graph and $G_{n} \in S_{n}$, by Proposition 2.1. Our next result shows that $S_{n}$ contains exponentially many planar graphs.

Theorem 2.1. For each $n \geq 11$ there are at least $2^{\left\lfloor\frac{n-11}{2}\right\rfloor} 2$-connected planar graphs in $S_{n}$.
Proof. In view of $G_{11}$ and $G_{12}$, we assume $n \geq 13$. Take $G_{n}$ and relabel the vertices of the $u_{2}-u_{3}$ path that passes through $w_{1}$ by $u_{2}=a_{0}, a_{1}, a_{2}, \ldots, a_{\ell-1}=w_{1}, a_{\ell}=u_{3}$; here $\ell=n-7$ if $n$ is odd and $\ell=n-8$ if $n$ is even. Note that $\ell$ is even. Let $E_{n}=\left\{a_{1} a_{\ell-2}, a_{2} a_{\ell-3}, \ldots, a_{\ell / 2-2} a_{\ell / 2+1}\right\}$. For every $E^{\prime} \subseteq E_{n}$, denote by $H_{n, E^{\prime}}$ the graph obtained from $G_{n}$ by adding the edges of $E^{\prime}$. Obviously $H_{n, E^{\prime}}$ is planar, $\Delta\left(H_{n, E^{\prime}}\right)=4$ and $\chi\left(H_{n, E^{\prime}}\right)=3$. Moreover, $H_{n, E^{\prime}}-\{x, y\}$ is bipartite if $\{x, y\}=\left\{v_{i}, v_{j}\right\}$ where $1 \leq i<j \leq 3$, which implies that $\operatorname{vs}_{\chi}\left(H_{n, E^{\prime}}\right)=2$. Also, $H_{n, E^{\prime}}-\left\{u_{2}, v_{2}, w_{2}\right\}$ is bipartite which gives ivs $\left(H_{n, E^{\prime}}\right) \leq 3$. On the other hand, since $G_{n}$ is a subgraph of $H_{n, E^{\prime}}$ of certain chromaticity and ivs $\chi_{\chi}\left(G_{n}\right)=3$, we have ivs $\chi\left(H_{n, E^{\prime}}\right)=3$ as well. Thus $H_{n, E^{\prime}} \in S_{n}$.


Figure 3: The graph $H_{18, E_{18}}$.
Let $E^{\prime}, E^{*} \subseteq E_{n}$, where $E^{\prime} \neq E^{*}$. We show that the graphs $H_{n, E^{\prime}}$ and $H_{n, E^{*}}$ are not isomorphic. This is obvious if $\left|E^{\prime}\right| \neq\left|E^{*}\right|$. So assume that $\left|E^{\prime}\right|=\left|E^{*}\right| \geq 1$. We show that Aut $\left(H_{n, E^{\prime}}\right)$, the group of automorphisms of $H_{n, E^{\prime}}$, (and also $\left.\operatorname{Aut}\left(H_{n, E^{*}}\right)\right)$ is trivial. That is, every automorphism of $H_{n, E^{\prime}}$ fixes all the vertices of $H_{n, E^{\prime}}$.

There are exactly 6 vertices of degree 4 in $H_{n, E^{\prime}}$, namely $u_{1}, u_{2}, u_{3}, v_{1}, v_{2}, v_{3}$. Since each of $u_{1}, u_{2}, u_{3}$ is in only one triangle in $H_{n, E^{\prime}}$ whereas each of $v_{1}, v_{2}, v_{3}$ is in three such triangles, every automorphism must preserve the sets $\left\{u_{1}, u_{2}, u_{3}\right\}$ and $\left\{v_{1}, v_{2}, v_{3}\right\}$. The vertices $u_{1}$ and $u_{2}$ are both adjacent to the same vertex of $H_{n, E^{\prime}}-\left\{v_{1}, v_{2}, v_{3}\right\}$. Also the vertices $u_{1}$ and $u_{3}$ are both adjacent to the same vertex of $H_{n, E^{\prime}}-\left\{v_{1}, v_{2}, v_{3}\right\}$. But $u_{2}$ and $u_{3}$ are not adjacent to the same vertex of $H_{n, E^{\prime}}-\left\{v_{1}, v_{2}, v_{3}\right\}$, because $n \geq 13$. Consequently, every automorphism of $H_{n, E^{\prime}}$ fixes $u_{1}$.

In view of $a_{\ell-1}\left(=w_{1}\right)$, the vertex $u_{3}$ has a neighbour of degree 2 which is not connected to $u_{1}$. If $u_{2}$ does not have such a neighbour, then every automorphism of $H_{n, E^{\prime}}$ fixes also $u_{2}$ and $u_{3}$. So assume that also $u_{2}$ has a neighbour of degree 2 which is not connected to $u_{1}$. Now start at $u_{2}$, proceed with the above mentioned neighbour of $u_{2}$ and construct a longest path $P_{2}$, interior vertices of which have all degree 2 . Analogously start at $u_{3}$, proceed with the above mentioned neighbour of $u_{3}$ and construct a longest path $P_{3}$, interior vertices of which have all degree 2 . Finally, let $i$ be the smallest index such that $a_{i} a_{\ell-1-i} \in E^{\prime}$. Then $P_{2}$ has length $i$ while $P_{3}$ has length $i+1$. Hence, every automorphism of $H_{n, E^{\prime}}$ must fix also $u_{2}$ and $u_{3}$. Consequently, every automorphism of $H_{n, E^{\prime}}$ fixes all the vertices of $H_{n, E^{\prime}}$, and so $H_{n, E^{\prime}}$ and $H_{n, E^{*}}$ are not isomorphic graphs.

Since $E_{n}$ has $\frac{\ell}{2}-2=\left\lfloor\frac{n-7}{2}\right\rfloor-2=\left\lfloor\frac{n-11}{2}\right\rfloor$ edges and every subset gives different graph, there are exactly $2^{\left\lfloor\frac{n-11}{2}\right\rfloor}$ nonisomorphic graphs $H_{n, E^{\prime}}$.

Remark. Observe that considering subsets of a zig-zag path $a_{l-1}, a_{1}, a_{l-2}, a_{2}, \ldots$ instead of set of isolated edges, one can obtain $2^{n-11}$ graphs satisfying the assumptions of Theorem 2.1, since the obtained graphs $G$ have $|\operatorname{Aut}(G)| \leq 2$, and still different subsets distinguish them.

We conclude the paper by presenting another, more general, construction of graphs $G$ with $\Delta(G)=4, \chi(G)=3$, $\operatorname{ivs}_{\chi}(G)=3$ and $\operatorname{vs}_{\chi}(G)=2$. Let $H$ be a bipartite graph with $\Delta(H) \leq 4$ such that there exists a cycle $C_{2 k} \subseteq H$ with $k \geq 3$ and a pair $a, b \in V\left(C_{2 k}\right)$ of non-adjacent vertices in $H$ on odd distance $d_{H}(a, b)$ and having $\operatorname{deg}_{H}(a)=\operatorname{deg}_{H}(b)=2$. Take the union of $H$ with a disjoint triangle $u v w$ and add the edges $a v, b v, a w, b w$. Denote the resulting graph by $G$ (see Figure 4).


Figure 4: The graph $G$ if $H=C_{6}$.

Proposition 2.2. If $H$ is of order $m$ then $G \in S_{m+3}$. Moreover, if $H$ is 2-connected (resp. planar) then $G$ has the same property.

Proof. Clearly, the order of $G$ is $m+3$. As $\Delta(H) \leq 4$ and $\operatorname{deg}_{H}(a)=\operatorname{deg}_{H}(b)=2$ we have $\Delta(G)=4$. In view of the triangle $u v w$, the graph $G$ is not bipartite. In order to show $\chi(G)=3$, note that $G-\{u, v, w\}$ is bipartite. Take a proper 2 -colouring $\varphi$ of $G-\{u, v, w\}$ with colours 1 and 3 such that, without loss of generality, $\varphi(a)=1$ and $\varphi(b)=3$ (here we use that $d(a, b)$ is odd). Now change the colour of $b$ to 1 and the colour of every $c \in N_{G-\{u, v, w\}}(b)$ to 2 . Note that by assigning the colour 1 to $u$, the colour 2 to $v$ and the colour 3 to $w$ we obtain a proper 3-colouring of $G$.

Let us show next that $\operatorname{ivs}_{\chi}(G)=3$. Since $G-\{a, b, u\}$ is bipartite, we have ivs $\chi_{\chi}(G) \leq 3$. Suppose there are non-adjacent vertices $x, y$ such that $G-\{x, y\}$ is bipartite. In view of the triangle $u v w$, the intersection $\{x, y\} \cap\{u, v, w\}$ is a singleton. We argue that this intersection is not the vertex $u$ due to the triangles $a v w$ and $b v w$. Let $P$ and $Q$ be the two $a-b$ paths in $C_{2 k}$, and recall that both these paths are of odd lengths. Consequently, each of the cycles $C^{\prime}=P \cup a v b, C^{\prime \prime}=P \cup a w b$, $C^{\prime \prime \prime}=Q \cup a v b$, and $C^{\prime \prime \prime \prime}=Q \cup a w b$ is odd. Hence $\{x, y\} \cap\{v, w\} \neq \emptyset$, which further implies that $\{x, y\} \cap\{a, b\}=\emptyset$. However, then at least one of the cycles $C^{\prime}, C^{\prime \prime}, C^{\prime \prime \prime}, C^{\prime \prime \prime \prime}$ appears in $G-\{x, y\}$. The obtained contradiction shows ivs $(G)=3$.

Finally, we prove that $\operatorname{vs}_{\chi}(G)=2$. Clearly $\operatorname{vs}_{\chi}(G) \leq 2$, because $G-\{v, w\}$ is bipartite. And since $\mathrm{vs}_{\chi}(G)=1$ implies $\operatorname{ivs}_{\chi}(G)=1$, we have $\operatorname{vs}_{\chi}(G)=2$.
Remark. Note in passing that the order of the bound $\left|S_{n}\right| \geq 2^{\left\lfloor\frac{n-11}{2}\right\rfloor}$ obtained in Theorem 2.1 is not (asymptotically) optimal. Propositions 2.1 and 2.2 enable one to construct connected planar graphs within $S_{n}$ with considerable ease. However, establishing a more precise asymptotic estimate of $\left|S_{n}\right|$ was not the focus of this short article; instead, the aim was simply to point out to the existence of exponentially many graphs in $S_{n}$.

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