# Research Article On chromatic vertex stability of 3-chromatic graphs with maximum degree 4

Martin Knor<sup>1</sup>, Mirko Petruševski<sup>2,\*</sup>, Riste Škrekovski<sup>3,4</sup>

<sup>1</sup>Slovak University of Technology in Bratislava, Bratislava, Slovakia

<sup>2</sup>Faculty of Mechanical Engineering, University Ss. Cyril and Methodius in Skopje, Skopje, Macedonia

<sup>3</sup>Faculty of Mathematics and Physics, University of Ljubljana, Ljubljana, Slovenia

<sup>4</sup>Faculty of Information Studies, Novo Mesto, Slovenia

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#### Abstract

The (independent) chromatic vertex stability  $(ivs_{\chi}(G)) vs_{\chi}(G)$  is the minimum size of (independent) set  $S \subseteq V(G)$  such that  $\chi(G-S) = \chi(G) - 1$ . The question of how large must the chromatic number  $\chi(G)$  of a graph G be, in terms of the maximum degree  $\Delta(G)$ , to ensure the equality  $ivs_{\chi}(G) = vs_{\chi}(G)$  was raised by Akbari et al. [*European J. Combin.* 102 (2022) #103504]; the authors showed that  $ivs_{\chi}(G) = vs_{\chi}(G)$  if  $\chi(G) \in {\Delta(G), \Delta(G) + 1}$ , and also pointed out to graphs with  $\chi(G) \leq (\Delta(G) + 1)/2$  for which  $ivs_{\chi}(G) > vs_{\chi}(G)$ . In the light of their findings, they raised the following problem: Is it true that  $\chi(G) \geq \Delta(G)/2 + 1$  always implies  $ivs_{\chi}(G) = vs_{\chi}(G)$ ? This threshold question was recently answered in the negative by Cambrie et al. [*arXiv*: 2203.13833v1, (2022)]. In this paper, we show that the smallest instance for counterexamples is the case  $(\chi(G), \Delta(G)) = (3, 4)$ , with the smallest possible order being 9 (and there are 30 such graphs). We construct exponentially many graphs G having  $\Delta(G) = 4$ ,  $\chi(G) = 3$ ,  $ivs_{\chi}(G) = 3$ , and  $vs_{\chi}(G) = 2$ .

Keywords: chromatic vertex stability; independent chromatic vertex stability; chromatic number; maximum degree.

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#### 1. Introduction

Let *G* be a graph. Its edge stability number,  $es_{\chi}(G)$ , is the minimum number of edges whose deletion results in a graph *H* with  $\chi(H) = \chi(G) - 1$ . The edge stability number was introduced in 1980 by Staton [9], and rediscovered in 2008 by Arumugam, Hamid, and Muthukamatchi [4]. For recent results on this invariant, see e.g. [1,3,6,8].

General concept of stability number appeared in [5], but the first paper on chromatic vertex stability number was written by Akbari, Beikmohammadi, Klavžar, and Movarraei in 2021, see [2]. The chromatic vertex stability  $vs_{\chi}(G)$  of G is the minimum number of vertices of G such that their deletion results in a graph H with  $\chi(H) = \chi(G) - 1$ . Analogously, the independent chromatic vertex stability  $vs_{\chi}(G)$  of G is the minimum number of independent vertices of G such that their deletion results in a graph H with  $\chi(H) = \chi(G) - 1$ . Obviously,  $vs_{\chi}(G) \leq ivs_{\chi}(G)$ . The main result of [2] is the following.

**Theorem 1.1.** If G is a graph with  $\chi(G) \in {\Delta(G), \Delta(G) + 1}$  then  $vs_{\chi}(G) = ivs_{\chi}(G)$ .

The authors defined the threshold function  $f(\Delta)$  as the smallest quantity such that, for any graph G of maximum degree  $\Delta$ , it must hold that  $vs_{\chi}(G) = ivs_{\chi}(G)$  provided that  $\chi(G) \ge f(\Delta)$ . Notice that Theorem 1.1 asserts  $f(\Delta) \le \Delta$ . They also showed that as soon as  $\chi(G) \le (\Delta(G) + 1)/2$  the equality  $vs_{\chi}(G) = ivs_{\chi}(G)$  need no longer be true, and consequently asked the following question (see Problem 3.2 in [2]).

**Problem 1.1.** Is it true that  $f(\Delta) \leq \frac{\Delta(G)}{2} + 1$ , that is, if G is a graph with  $\chi(G) \geq \frac{\Delta(G)}{2} + 1$  does it then always hold that  $vs_{\chi}(G) = ivs_{\chi}(G)$ ?

The question of Problem 1.1 was recently answered in the negative by Cambie et al. [7], who proved that  $f(\Delta) = \Delta$  for  $3 \le \Delta \le 10$ . Additionally, they determined the threshold  $f(\Delta)$  to within two values (and indeed sometimes a unique value) for graphs of sufficiently large maximum degree.

In this paper, we focus on the smallest value  $\chi(G) = 3$  for which counterexamples to Problem 1.1 exist, i.e. we consider the case when  $\chi(G) = 3$  and  $\Delta(G) = 4$ . A simple 'ladder-like' counterexample on 9 vertices is depicted in Figure 1. Notice that the ladder part of the counterexample can be of any length 4k for  $k \ge 2$ .

<sup>\*</sup>Corresponding author (mirko.petrushevski@gmail.com).



Figure 1: A graph G with  $\Delta(G) = 4$  and  $\chi(G) = 3$ . The blue independent 3-set realizes  $ivs_{\chi}(G) = 3$  and the yellow 2-set realizes  $vs_{\chi}(G) = 2$ .

Our main result is the following.

**Theorem 1.2.** For each  $n \ge 9$  there are at least  $\max\left\{1, 2^{\lfloor \frac{n-11}{2} \rfloor}\right\}$  planar graphs with  $\chi(G) = 3$ ,  $\Delta(G) = 4$ ,  $\operatorname{ivs}_{\chi}(G) = 3$  and  $\operatorname{vs}_{\chi}(G) = 2$ .

Observe that if  $\chi(G) = 3$  and  $\Delta(G) = 4$ , then  $\chi(G) = \frac{\Delta(G)}{2} + 1 = \Delta(G) - 1$ . Hence, the bound on  $\chi(G)$  in Theorem 1.1 cannot be relaxed when  $\chi(G) = 3$ .

#### 2. Proofs

We start with a pair of simple observations followed by a couple of lemmas.

**Observation 2.1.** For every graph G,  $ivs_{\chi}(G)$  equals the minimum size of a colour class over all proper  $\chi(G)$ -colourings of G. Hence  $|V(G)| \ge ivs_{\chi}(G) \cdot \chi(G)$ .

*Proof.* Let us first notice that there exists a proper  $\chi(G)$ -colouring of G with a colour class of size  $ivs_{\chi}(G)$ . Indeed, take  $S \subseteq V(G)$  to be an independent set of size  $|S| = ivs_{\chi}(G)$  such that  $\chi(G - S) = \chi(G) - 1$ . Use a proper  $(\chi(G) - 1)$ -colouring of G - S and assign to all vertices of S a new colour. So  $ivs_{\chi}(G)$  is not less than the minimum size of a colour class over all proper  $\chi(G)$ -colourings of G.

Contrarily, consider a proper  $\chi(G)$ -colouring of G which minimizes the size of a colour class, and let S be such a minimum colour class. Then S is an independent subset of V(G) and  $\chi(G-S) \leq \chi(G) - 1$ . In fact, we must have equality here for otherwise G would admit a proper  $(\chi(G) - 1)$ -colouring. So  $ivs_{\chi}(G)$  is also not more than the minimum size of a colour class over all proper  $\chi(G)$ -colourings of G, which proves our point.

The inequality  $|V(G)| \ge ivs_{\chi}(G) \cdot \chi(G)$  is now an immediate consequence.

**Observation 2.2.** If  $\Delta(G) \leq 2$  then  $vs_{\chi}(G) = ivs_{\chi}(G)$ .

*Proof.* We may assume that  $vs_{\chi}(G) \ge 2$ . Indeed, if  $vs_{\chi}(G) = 1$  then obviously  $ivs_{\chi}(G) = 1$  as well. We may also assume that *G* is connected. Then *G* is either a path or an even cycle. In either case

$$s_{\chi}(G) = ivs_{\chi}(G) = \left\lfloor \frac{|V(G)|}{2} \right\rfloor.$$

As already mentioned, we are interested in finding graphs G for which  $\chi(G) \ge \frac{\Delta(G)}{2} + 1$  and  $ivs_{\chi}(G) > vs_{\chi}(G)$ . Our first lemma establishes some implications for the order and the considered stability parameters.

**Lemma 2.1.** If  $\operatorname{ivs}_{\chi}(G) > \operatorname{vs}_{\chi}(G)$  and  $\chi(G) \ge \frac{\Delta(G)}{2} + 1$  then  $|V(G)| \ge 9$ ,  $\operatorname{ivs}_{\chi}(G) \ge 3$ ,  $\operatorname{vs}_{\chi}(G) \ge 2$  and  $\chi(G) \ge 3$ . Moreover, if |V(G)| = 9 then  $\operatorname{ivs}_{\chi}(G) = 3$ ,  $\operatorname{vs}_{\chi}(G) = 2$  and  $\chi(G) = 3$ .

*Proof.* Since  $ivs_{\chi}(G) > vs_{\chi}(G)$ , we must have  $vs_{\chi}(G) \ge 2$  and consequently  $ivs_{\chi}(G) \ge 3$ .

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If  $\chi(G) \leq 2$  then from  $\chi(G) \geq \frac{\Delta(G)}{2} + 1$  we get  $\Delta(G) \leq 2(\chi(G) - 1) \leq 2$ , which in view of Observation 2.2 contradicts  $ivs_{\chi}(G) > vs_{\chi}(G)$ . Hence  $\chi(G) \geq 3$ .

From the inequality stated in Observation 2.1, it follows that  $|V(G)| \ge ivs_{\chi}(G) \cdot \chi(G) \ge 3 \cdot 3$ , that is,  $|V(G)| \ge 9$ . And if |V(G)| = 9 then  $\chi(G) = ivs_{\chi}(G) = 3$  and  $vs_{\chi}(G) = 2$ .

 $\square$ 

Figure 2 depicts two graphs, respectively denoted by  $G_9$  and  $G_{10}$  in regard to their orders. The former one can be obtained from the octahedron graph by subdividing the edges of a triangle. It has  $\Delta(G_9) = 4$ ,  $\chi(G_9) = \operatorname{ivs}_{\chi}(G_9) = 3$  and  $\operatorname{vs}_{\chi}(G_9) = 2$ . Observe that  $\chi(G_9 - \{x, y\}) = 2$  if and only if  $\{x, y\} = \{v_i, v_j\}$ , where  $1 \le i < j \le 3$ , and for an independent set of vertices  $\{x, y, z\}$  we have  $\chi(G_9 - \{x, y, z\}) = 2$  if and only if  $\{x, y, z\} = \{u_i, v_i, w_i\}$ , where  $1 \le i \le 3$ . The graph  $G_{10}$  is obtained from  $G_9$  by adding the vertex q and connecting it to  $w_2$  and  $w_3$ . It also has  $\Delta(G_{10}) = 4$ ,  $\chi(G_{10}) = \operatorname{ivs}_{\chi}(G_{10}) = 3$  and  $\operatorname{vs}_{\chi}(G_{10}) = 2$ . Again  $\chi(G_{10} - \{x, y, z\}) = 2$  if and only if  $\{x, y\} = \{v_i, v_j\}$ , where  $1 \le i < j \le 3$ , and for an independent set of vertices  $\{x, y, z\}$  we have  $\chi(G_{10} - \{x, y, z\}) = 2$  if and only if  $\{x, y, z\} = \{u_i, v_i, w_i\}$  where  $2 \le i \le 3$ .



Figure 2: The graph  $G_9$  (left) and the graph  $G_{10}$  (right).

Our second lemma concerns the case  $\chi(G) = ivs_{\chi}(G) = 3$  and  $vs_{\chi}(G) = 2$ . Under the assumption  $\chi(G) \ge \frac{\Delta(G)}{2} + 1$  we establish the maximum degree of G and the vertex degrees of every 2-set which realizes  $vs_{\chi}(G)$ .

**Lemma 2.2.** Let G be a graph with  $\chi(G) = ivs_{\chi}(G) = 3$ ,  $vs_{\chi}(G) = 2$  and  $\chi(G) \ge \frac{\Delta(G)}{2} + 1$ . Then  $\Delta(G) = 4$  and for every  $v_1, v_2 \in V(G)$  such that  $\chi(G - \{v_1, v_2\}) = 2$  we have  $\deg_G(v_1) = \deg_G(v_2) = 4$ .

*Proof.* Let G satisfy the assumptions of the lemma and let  $v_1, v_2 \in V(G)$  be such that  $\chi(G - \{v_1, v_2\}) = 2$ . Then

$$\Delta(G) \le 2(\chi(G) - 1) = 4.$$

For argument's sake, suppose that  $\deg_G(v_1) \leq 3$ . Since  $\operatorname{ivs}_{\chi}(G) = 3 > 2$ , we have  $v_1v_2 \in E(G)$ . Further, since  $\chi(G - v_2) = 3$ , there must be an odd cycle in  $G - v_2$ ; moreover, every such cycle passes through  $v_1$  since  $G - \{v_1, v_2\}$  is bipartite. Consequently  $\deg_G(v_1) = 3$ .

Let  $u_1, u_2$  be the neighbours of  $v_1$  in  $G - v_2$ . Since every odd cycle in  $G - v_2$  passes through both  $u_1, u_2$ , we conclude that  $G - \{u_i, v_2\}$  is bipartite as well,  $i \in \{1, 2\}$ . From  $ivs_{\chi}(G) = 3 > 2$  it follows that  $u_1v_2, u_2v_2 \in E(G)$ . Moreover,  $u_1u_2 \notin E(G)$ , for otherwise  $v_1, v_2, u_1, u_2$  induces a  $K_4$ , implying  $\chi(G) \ge 4$ .

Since  $ivs_{\chi}(G) = 3$ , there must be an odd cycle in  $G - \{u_1, u_2\}$ . This cycle cannot pass through  $v_1$  since  $deg_{G-\{u_1, u_2\}}(v_1) = 1$ . If this cycle does not pass through  $v_2$  as well then it is in  $G - \{v_1, v_2\}$  which means that  $\chi(G - \{v_1, v_2\}) = 3$ , a contradiction. Hence, there is an odd cycle passing through  $v_2$  in  $G - \{u_1, u_2, v_1\}$ , which means that  $deg_G(v_2) \ge 5$ . This contradiction settles the lemma.

A computer search shows that there are precisely 30 graphs G of order 9 and having  $\Delta(G) = 4$ ,  $\chi(G) = 3$ ,  $ivs_{\chi}(G) = 3$ and  $vs_{\chi}(G) = 2$ . Several of them (including  $G_9$ ) are planar and four are obtained by adding an edge to another graph from the same collection.

For every  $n \ge 9$  let  $S_n$  be the set of graphs G on n vertices such that  $\Delta(G) = 4$ ,  $\chi(G) = 3$ ,  $\operatorname{ivs}_{\chi}(G) = 3$  and  $\operatorname{vs}_{\chi}(G) = 2$ . Thus  $G_9 \in S_9$  and  $G_{10} \in S_{10}$ . By  $C_{\chi}(G)$  we denote the set of vertices  $x \in V(G)$  such that there is some  $y \in V(G)$  for which  $\chi(G - \{x, y\}) = 2$ ; note that every such y is a neighbour of x. For example,  $C_{\chi}(G_9) = C_{\chi}(G_{10}) = \{v_1, v_2, v_3\}$ . In view of our next result, for every  $n \ge 9$  there is a planar graph in  $S_n$  which is topologically equivalent to  $G_9$  or  $G_{10}$ . **Proposition 2.1.** Let  $G \in S_n$  and  $e_1, e_2, \ldots, e_t \in E(G) \setminus E([C_{\chi}(G)])$ , i.e., each  $e_i$  has at most one endvertex in  $C_{\chi}(G)$ . Let  $n_1, n_2, \ldots, n_t$  be positive even integers. For every  $i, 1 \le i \le t$ , subdivide  $e_i$  with  $n_i$  new vertices, and denote the resulting graph by H. Then  $\Delta(H) = 4$ ,  $\chi(H) = 3$ ,  $ivs_{\chi}(H) = 3$  and  $vs_{\chi}(H) = 2$ . In other words,  $H \in S_{n+(n_1+\cdots+n_t)}$ .

*Proof.* Obviously  $\Delta(H) = \Delta(G) = 4$ . Since  $\chi(G) = 3$ , the graph G has an odd cycle. Since to any edge of this cycle we added an even number (possibly zero) of vertices, H also has an odd cycle; thus  $\chi(H) \ge 3$ . Moreover, if  $S \subseteq V(G)$  is such that G - S is bipartite then H - S is bipartite as well. Hence,  $\chi(H) = 3$ ,  $\operatorname{ivs}_{\chi}(H) \le 3$  and  $\operatorname{vs}_{\chi}(H) \le 2$ .

If there is some  $v \in V(H)$  such that  $\chi(H - v) = 2$ , then  $v \notin V(G)$ . So v is obtained by subdividing an edge, say xy, of G. However, as H - v is bipartite, both G - x and G - y are bipartite, a contradiction. Hence,  $vs_{\chi}(H) = 2$ .

Finally, let us show that  $ivs_{\chi}(H) = 3$ . Supposing the opposite, there are  $u, v \in V(H)$  such that  $\chi(H - \{u, v\}) = 2$  and  $uv \notin E(H)$ . It cannot be that both u and v are in V(G), because we did not subdivide edges connecting vertices of  $C_{\chi}(G)$ . So we may assume that v is obtained by subdividing an edge xy of G, where  $y \notin C_{\chi}(G)$ . Since  $H - \{u, v\}$  is bipartite, so is  $H - \{u, y\}$ . But then u cannot be a vertex of G as well. Hence, u is obtained by subdividing an edge wz of G. As  $H - \{u, y\}$  is bipartite, so is  $H - \{z, y\}$ . However, this contradicts the fact that  $y \notin C_{\chi}(G)$ .

From Proposition 2.1 we deduce that  $S_n \neq \emptyset$  for every  $n \ge 9$ . Indeed, if n is odd then take  $G_9$ , subdivide the edge  $u_2w_1$  with n - 9 new vertices and denote the resulting graph by  $G_n$ . Analogously if n is even then take  $G_{10}$ , subdivide the edge  $u_2w_1$  with n - 10 new vertices and denote the resulting graph by  $G_n$ . Then  $G_n$  is a connected planar graph and  $G_n \in S_n$ , by Proposition 2.1. Our next result shows that  $S_n$  contains exponentially many planar graphs.

## **Theorem 2.1.** For each $n \ge 11$ there are at least $2^{\lfloor \frac{n-11}{2} \rfloor}$ 2-connected planar graphs in $S_n$ .

*Proof.* In view of  $G_{11}$  and  $G_{12}$ , we assume  $n \ge 13$ . Take  $G_n$  and relabel the vertices of the  $u_2 - u_3$  path that passes through  $w_1$  by  $u_2 = a_0, a_1, a_2, \ldots, a_{\ell-1} = w_1, a_\ell = u_3$ ; here  $\ell = n - 7$  if n is odd and  $\ell = n - 8$  if n is even. Note that  $\ell$  is even. Let  $E_n = \{a_1 a_{\ell-2}, a_2 a_{\ell-3}, \ldots, a_{\ell/2-2} a_{\ell/2+1}\}$ . For every  $E' \subseteq E_n$ , denote by  $H_{n,E'}$  the graph obtained from  $G_n$  by adding the edges of E'. Obviously  $H_{n,E'}$  is planar,  $\Delta(H_{n,E'}) = 4$  and  $\chi(H_{n,E'}) = 3$ . Moreover,  $H_{n,E'} - \{x, y\}$  is bipartite if  $\{x, y\} = \{v_i, v_j\}$  where  $1 \le i < j \le 3$ , which implies that  $v_{s_{\chi}}(H_{n,E'}) = 2$ . Also,  $H_{n,E'} - \{u_2, v_2, w_2\}$  is bipartite which gives  $iv_{s_{\chi}}(H_{n,E'}) \le 3$ . On the other hand, since  $G_n$  is a subgraph of  $H_{n,E'}$  of certain chromaticity and  $iv_{s_{\chi}}(G_n) = 3$ , we have  $iv_{s_{\chi}}(H_{n,E'}) = 3$  as well. Thus  $H_{n,E'} \in S_n$ .



Figure 3: The graph  $H_{18,E_{18}}$ .

Let  $E', E^* \subseteq E_n$ , where  $E' \neq E^*$ . We show that the graphs  $H_{n,E'}$  and  $H_{n,E^*}$  are not isomorphic. This is obvious if  $|E'| \neq |E^*|$ . So assume that  $|E'| = |E^*| \ge 1$ . We show that  $\operatorname{Aut}(H_{n,E'})$ , the group of automorphisms of  $H_{n,E'}$ , (and also  $\operatorname{Aut}(H_{n,E^*})$ ) is trivial. That is, every automorphism of  $H_{n,E'}$  fixes all the vertices of  $H_{n,E'}$ .

There are exactly 6 vertices of degree 4 in  $H_{n,E'}$ , namely  $u_1, u_2, u_3, v_1, v_2, v_3$ . Since each of  $u_1, u_2, u_3$  is in only one triangle in  $H_{n,E'}$  whereas each of  $v_1, v_2, v_3$  is in three such triangles, every automorphism must preserve the sets  $\{u_1, u_2, u_3\}$  and  $\{v_1, v_2, v_3\}$ . The vertices  $u_1$  and  $u_2$  are both adjacent to the same vertex of  $H_{n,E'} - \{v_1, v_2, v_3\}$ . Also the vertices  $u_1$  and  $u_3$  are both adjacent to the same vertex of  $H_{n,E'} - \{v_1, v_2, v_3\}$ . But  $u_2$  and  $u_3$  are not adjacent to the same vertex of  $H_{n,E'} - \{v_1, v_2, v_3\}$ , because  $n \ge 13$ . Consequently, every automorphism of  $H_{n,E'}$  fixes  $u_1$ . In view of  $a_{\ell-1}(=w_1)$ , the vertex  $u_3$  has a neighbour of degree 2 which is not connected to  $u_1$ . If  $u_2$  does not have such a neighbour, then every automorphism of  $H_{n,E'}$  fixes also  $u_2$  and  $u_3$ . So assume that also  $u_2$  has a neighbour of degree 2 which is not connected to  $u_1$ . Now start at  $u_2$ , proceed with the above mentioned neighbour of  $u_2$  and construct a longest path  $P_2$ , interior vertices of which have all degree 2. Analogously start at  $u_3$ , proceed with the above mentioned neighbour of  $u_3$  and construct a longest path  $P_3$ , interior vertices of which have all degree 2. Finally, let *i* be the smallest index such that  $a_i a_{\ell-1-i} \in E'$ . Then  $P_2$  has length *i* while  $P_3$  has length i+1. Hence, every automorphism of  $H_{n,E'}$  must fix also  $u_2$  and  $u_3$ . Consequently, every automorphism of  $H_{n,E'}$  fixes all the vertices of  $H_{n,E'}$ , and so  $H_{n,E'}$  and  $H_{n,E^*}$  are not isomorphic graphs.

Since  $E_n$  has  $\frac{\ell}{2} - 2 = \lfloor \frac{n-7}{2} \rfloor - 2 = \lfloor \frac{n-11}{2} \rfloor$  edges and every subset gives different graph, there are exactly  $2^{\lfloor \frac{n-11}{2} \rfloor}$  nonisomorphic graphs  $H_{n,E'}$ .

**Remark.** Observe that considering subsets of a zig-zag path  $a_{l-1}, a_1, a_{l-2}, a_2, \ldots$  instead of set of isolated edges, one can obtain  $2^{n-11}$  graphs satisfying the assumptions of Theorem 2.1, since the obtained graphs G have  $|\operatorname{Aut}(G)| \leq 2$ , and still different subsets distinguish them.

We conclude the paper by presenting another, more general, construction of graphs G with  $\Delta(G) = 4$ ,  $\chi(G) = 3$ ,  $ivs_{\chi}(G) = 3$  and  $vs_{\chi}(G) = 2$ . Let H be a bipartite graph with  $\Delta(H) \leq 4$  such that there exists a cycle  $C_{2k} \subseteq H$  with  $k \geq 3$ and a pair  $a, b \in V(C_{2k})$  of non-adjacent vertices in H on odd distance  $d_H(a, b)$  and having  $\deg_H(a) = \deg_H(b) = 2$ . Take the union of H with a disjoint triangle uvw and add the edges av, bv, aw, bw. Denote the resulting graph by G (see Figure 4).



Figure 4: The graph G if  $H = C_6$ .

**Proposition 2.2.** If H is of order m then  $G \in S_{m+3}$ . Moreover, if H is 2-connected (resp. planar) then G has the same property.

*Proof.* Clearly, the order of G is m + 3. As  $\Delta(H) \leq 4$  and  $\deg_H(a) = \deg_H(b) = 2$  we have  $\Delta(G) = 4$ . In view of the triangle uvw, the graph G is not bipartite. In order to show  $\chi(G) = 3$ , note that  $G - \{u, v, w\}$  is bipartite. Take a proper 2-colouring  $\varphi$  of  $G - \{u, v, w\}$  with colours 1 and 3 such that, without loss of generality,  $\varphi(a) = 1$  and  $\varphi(b) = 3$  (here we use that d(a, b) is odd). Now change the colour of b to 1 and the colour of every  $c \in N_{G-\{u,v,w\}}(b)$  to 2. Note that by assigning the colour 1 to u, the colour 2 to v and the colour 3 to w we obtain a proper 3-colouring of G.

Let us show next that  $ivs_{\chi}(G) = 3$ . Since  $G - \{a, b, u\}$  is bipartite, we have  $ivs_{\chi}(G) \leq 3$ . Suppose there are non-adjacent vertices x, y such that  $G - \{x, y\}$  is bipartite. In view of the triangle uvw, the intersection  $\{x, y\} \cap \{u, v, w\}$  is a singleton. We argue that this intersection is not the vertex u due to the triangles avw and bvw. Let P and Q be the two a - b paths in  $C_{2k}$ , and recall that both these paths are of odd lengths. Consequently, each of the cycles  $C' = P \cup avb$ ,  $C'' = P \cup awb$ ,  $C''' = Q \cup avb$ , and  $C'''' = Q \cup awb$  is odd. Hence  $\{x, y\} \cap \{v, w\} \neq \emptyset$ , which further implies that  $\{x, y\} \cap \{a, b\} = \emptyset$ . However, then at least one of the cycles C', C'', C''', C'''' appears in  $G - \{x, y\}$ . The obtained contradiction shows  $ivs_{\chi}(G) = 3$ .

Finally, we prove that  $vs_{\chi}(G) = 2$ . Clearly  $vs_{\chi}(G) \leq 2$ , because  $G - \{v, w\}$  is bipartite. And since  $vs_{\chi}(G) = 1$  implies  $ivs_{\chi}(G) = 1$ , we have  $vs_{\chi}(G) = 2$ .

**Remark.** Note in passing that the order of the bound  $|S_n| \ge 2^{\lfloor \frac{n-11}{2} \rfloor}$  obtained in Theorem 2.1 is not (asymptotically) optimal. Propositions 2.1 and 2.2 enable one to construct connected planar graphs within  $S_n$  with considerable ease. However, establishing a more precise asymptotic estimate of  $|S_n|$  was not the focus of this short article; instead, the aim was simply to point out to the existence of exponentially many graphs in  $S_n$ .

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